

## Recap

→ **channel twirling**: averaging operation on channels that enforces covariance property

→ Let  $\mathcal{N}: A \rightarrow B$  be a channel,  $G$  a finite group with unitary

reps  $U_g$  on  $\mathcal{H}_A$ ,  $V_g$  on  $\mathcal{H}_B$ : 
$$\mathcal{N}_G(\cdot) = \frac{1}{|G|} \sum_{g \in G} V_g^\dagger \mathcal{N}(U_g \cdot U_g^\dagger) V_g$$

→  $\mathcal{N}_G$  is  $(G, U_g, V_g)$ -covariant.

→ For continuous groups: need suitable averaging operation

→ Relevant case for us: **compact Lie groups** (such as unitary group)

For such groups there exists a unique "uniform" probability measure

$\mu: \sigma_G \rightarrow \mathbb{R}^+$  satisfying left- and right-invariance:

$$\mu(S) = \mu(gS) = \mu(Sg) \quad \forall g \in G, S \in \sigma_G$$

→ This measure is called **Haar measure** and gives rise to the

**Haar integral** 
$$\int_G d\mu(g) f(g)$$

→ For finite groups,  $\mu$  is the counting measure and

$$\int_G d\mu(g) f(g) = \frac{1}{|G|} \sum_{g \in G} f(g)$$

→ Channel twirl: 
$$\mathcal{N}_G = \int_G d\mu(g) V_g^\dagger \mathcal{N}(U_g \cdot U_g^\dagger) V_g$$

→ Result:  $\mathcal{N}_{\text{dep}(A)} = \int_{U(A)} dU U^\dagger \mathcal{N}(U \cdot U^\dagger) U$  is a depolarizing channel.

### Lemma 22

An irrep  $(\varphi, R)$  of a group  $G$  (finite or compact)

forms a so-called 1-design:  $\frac{1}{|G|} \sum_{g \in G} \varphi(g) X \varphi(g)^\dagger = \frac{\text{tr} X}{d} \mathbb{1}_R$ ,

where  $X \in \mathcal{B}(R)$ , and  $d = \dim R$ .

Proof: Schur's lemma. □

### Prop 23

Let  $\mathcal{N}: \mathcal{B}(X) \rightarrow \mathcal{B}(K)$  be a  $(G, U_g, V_g)$ -covariant channel where  $g \mapsto V_g$  on  $K$  is irreducible.

Then  $\mathcal{N}$  satisfies  $\mathcal{N}\left(\frac{1}{|X|} \mathbb{1}_X\right) = \frac{1}{|K|} \mathbb{1}_K$ .

If furthermore  $|X| = |K|$  (i.e.,  $X \cong K$ ), then  $\mathcal{N}$  is unital.

Proof:  $V_g \mathcal{N}(\cdot) V_g^\dagger = \mathcal{N}(U_g \cdot U_g^\dagger) \quad \forall g \in G$

$$\mathcal{N}(\mathbb{1}_X) = \mathcal{N}(U_g \mathbb{1}_X U_g^\dagger) = V_g \mathcal{N}(\mathbb{1}_X) V_g^\dagger \quad \forall g \in G$$

Lemma 22  
 $\Rightarrow$   
 (Schur's lemma)

$$\mathcal{N}(\mathbb{1}_X) = \frac{\text{tr} \mathcal{N}(\mathbb{1}_X)}{|K|} \mathbb{1}_K = \frac{|X|}{|K|} \mathbb{1}_K$$

$\uparrow$   
 $\text{tr} \mathcal{N}(\mathbb{1}_X) = \text{tr} \mathbb{1}_X$

$$\Rightarrow \mathcal{N}\left(\frac{1}{|X|} \mathbb{1}_X\right) = \frac{1}{|K|} \mathbb{1}_K, \text{ and } \mathcal{N}(\mathbb{1}_X) = \mathbb{1}_K \text{ if } |X| = |K|. \quad \square$$

For irreducibly covariant channels, the Holevo information assumes a simple form.

Reminder: in general,

$$\chi(N) = \max_{\{p_x, \rho_x\}} \left\{ S(N(\sum_x p_x \rho_x)) - \sum_x p_x S(N(\rho_x)) \right\}$$

**Lemma 24**

The maximization in  $\chi(N)$  can be restricted to pure state ensembles.

Proof: Let  $\{p_x, \rho_x\}$  be a (mixed) state ensemble achieving  $\chi(N)$ :

$$\chi(N) = S(\sum_x p_x N(\rho_x)) - \sum_x p_x S(N(\rho_x))$$

Let  $\rho_{xA} = \sum_x p_x |x\rangle\langle x|_X \otimes \rho_A^x$  (so-called classical-quantum state)

$$= \begin{pmatrix} p_1 \rho_A^1 & & & \\ 0 & p_2 \rho_A^2 & & \\ & & \ddots & \\ & & & & \end{pmatrix} = \bigoplus_x p_x \rho_A^x$$

$$\sigma_{xB} = (\text{id}_X \otimes N)(\rho_{xA}) = \sum_x p_x |x\rangle\langle x|_X \otimes N(\rho_A^x)$$

$$S(N(\sum_x p_x \rho_A^x)) = S(\sigma_B) \quad - \sum_x p_x \log p_x = H(\{p_x\})$$

$$\sum_x p_x S(N(\rho_A^x)) = S(\sigma_{xB}) - S(\sigma_X)$$

$$S(\mathcal{N}(\sum_x p_x \mathcal{P}_A^x)) = S(\sigma_B) \quad - \sum_x p_x \log p_x = H(\{p_x\}) \quad \text{Shannon entropy}$$

$$\sum_x p_x S(\mathcal{N}(\mathcal{P}_A^x)) = S(\sigma_{XB}) - S(\sigma_X)$$

$$\gamma(N) = S(\sigma_B) - (S(\sigma_{XB}) - S(\sigma_X))$$

$$= S(\sigma_B) + S(\sigma_X) - S(\sigma_{XB})$$

$$= I(X; B)_\sigma \quad \text{mutual information (between } X \text{ and } B)$$

$$\forall x: \mathcal{P}_A^x = \sum_y p_{y|x} |y_{x,y}\rangle \langle y_{x,y}|$$

$$\sigma_{XB} = \sum_{x,y} p_x p_{y|x} |x\rangle \langle x| \otimes |y\rangle \langle y| \otimes \mathcal{N}(|y_{x,y}\rangle \langle y_{x,y}|)$$

$$\sigma_{XB} = \tau_y \sigma_{XB}$$

$$I(XY; B)_\sigma \geq I(X; B)_\sigma = \underline{\gamma(N)} = I(Y; B)$$

↑

data-processing inequality (will be proved and discussed

in Q(2))

□

**Prop 25** Let  $N: B(\mathcal{X}) \rightarrow B(\mathcal{K})$  be a  $(G, U_g, V_g)$ -covariant channel.

i) If  $g \mapsto U_g$  on  $\mathcal{X}$  is irreducible, then  $(d_1 = \dim \mathcal{X})$

$$\chi(N) = S\left(N\left(\frac{1}{d} \mathbb{1}_{\mathcal{X}}\right)\right) - \min_{|\mathcal{Y}\rangle} S(N(|\mathcal{Y}\rangle)).$$

ii) if in addition  $g \mapsto V_g$  is irred., then  $N\left(\frac{1}{d_1} \mathbb{1}_{\mathcal{X}}\right) = \frac{1}{d_2} \mathbb{1}_{\mathcal{K}}$  ( $d_2 = |\mathcal{K}|$ )

$$\chi(N) = \log_2 d_2 - \min_{|\mathcal{Y}\rangle} S(N(|\mathcal{Y}\rangle)).$$

Proof: i) Prop 24: let  $\{p_x, |\varphi_x\rangle\langle\varphi_x|\}$  achieve max. in  $\chi(N)$ :

$$\chi(N) = S\left(\sum_x p_x N(|\varphi_x\rangle)\right) - \sum_x p_x S(N(|\varphi_x\rangle))$$

we have the following bounds:

$$\begin{aligned} \cdot) \sum_x p_x S(N(|\varphi_x\rangle)) &\geq \min_{|\mathcal{Y}\rangle} S(N(|\mathcal{Y}\rangle)) \\ &\geq \min_{|\mathcal{Y}\rangle} S(N(|\mathcal{Y}\rangle)) \end{aligned}$$

$$S(\rho) = S(U \rho U^\dagger) \quad \forall \rho \geq 0$$

Unitary

$$\cdot) \forall g \quad S\left(\sum_x p_x N(|\varphi_x\rangle)\right) = S\left(\sum_x p_x V_g N(|\varphi_x\rangle) V_g^\dagger\right)$$

$$S\left(\sum_x p_x \mathcal{N}(\varphi_x)\right) = \frac{1}{|G|} \sum_{g \in G} S\left(\sum_x p_x \underbrace{V_g \mathcal{N}(\varphi_x) V_g^\dagger}_{\mathcal{N}(U_g \varphi_x U_g^\dagger)}\right)$$

$$= \frac{1}{|G|} \sum_{g \in G} S\left(\sum_x p_x \mathcal{N}(U_g \varphi_x U_g^\dagger)\right)$$

von Neumann entropy is concave (Ex. sheet 3)  $\rightarrow$

$$\leq S\left(\sum_x p_x \underbrace{\mathcal{N}\left(\frac{1}{|G|} \sum_{g \in G} U_g \varphi_x U_g^\dagger\right)}_{= \frac{1}{d} \mathbb{1}_X \text{ by lemma 22}}\right)$$

$$= S\left(\mathcal{N}\left(\frac{1}{d} \mathbb{1}_X\right)\right)$$

In summary,  $\chi(N) = S\left(\sum_x p_x \mathcal{N}(\varphi_x)\right) - \sum_x p_x S(\mathcal{N}(\varphi_x))$

$$\leq S\left(\mathcal{N}\left(\frac{1}{d} \mathbb{1}_X\right)\right) - \min_{|\varphi\rangle} S(\mathcal{N}(\varphi))$$

This value is achieved by the following pure-state ensemble:

Let  $|\varphi\rangle$  achieve the minimum in  $\min_{|\varphi\rangle} S(\mathcal{N}(\varphi))$

$$S(\mathcal{N}(\varphi))$$

Define  $|\varphi_g\rangle = U_g |\varphi\rangle$  and  $p_g = \frac{1}{|G|}$

$$S(V_g \mathcal{N}(\varphi) V_g^\dagger)$$

$$\chi(\{p_g, |\varphi_g\rangle\}, N) = \underbrace{S\left(\sum_g \frac{1}{|G|} \mathcal{N}(U_g \varphi U_g^\dagger)\right)}_{\mathcal{N}\left(\frac{1}{d} \mathbb{1}\right)} - \frac{1}{|G|} \sum_{g \in G} S(\mathcal{N}(U_g \varphi U_g^\dagger))$$

ii) If  $g \mapsto V_g$  is irreducible as well, then by Prop 23,

$$\mathcal{N}\left(\frac{1}{d_1} \mathbb{1}_X\right) = \frac{1}{d_2} \mathbb{1}_U, \quad (d_1 = |X|, d_2 = |U|)$$

so that  $S\left(\mathcal{N}\left(\frac{1}{d_1} \mathbb{1}_X\right)\right) = S\left(\frac{1}{d_2} \mathbb{1}_U\right) = \log d_2.$  □