

Recap

- Entanglement-breaking channel: $(\text{id}_R \otimes N)(\rho_{RA}) \in \text{SEP}$ for all ρ_{RA}
- N EB \Leftrightarrow Choi op τ^N is separable
 - \Leftrightarrow there is a trans. map. $\{K_i\}$ with $\text{rk } K_i = 1 \ \forall i$
 - $\Leftrightarrow N$ is a measure-and-prepare channel:
 \exists POVM $\{E_i\}$ and states $\{\sigma_i\}$ s.t. $N(\rho) = \sum_i \text{tr}(E_i \rho) \sigma_i$
- EB channels can never generate entanglement \Rightarrow zero quantum cap.
- EB channels break entanglement between different inputs
 - \Rightarrow classical capacity $C(N) = \chi(N)$ Holevo information
- However: computing $\chi(N)$ is NP-hard :-)
- In general, deciding separability is NP-hard
 - \Rightarrow need for easily computable separability criteria
- Positive partial transpose (PPT): $\rho_{AB}^T \geq 0$
 $\rho_{AB} \text{ SEP} \Rightarrow \rho_{AB} \text{ PPT}$
- If $|A| \cdot |B| \leq 6$, also sufficient criterion: $\rho_{AB} \text{ PPT} \Rightarrow \rho_{AB} \text{ SEP}$
- PPT-channels: $(\text{id}_R \otimes N)(\rho_{RA})$ PPT for all ρ_{RA}
- N EB $\Rightarrow N$ PPT

Prop 12 TFAE:

a) \mathcal{N} is PPT

b) Choi op $\tau^{\mathcal{N}}$ is PPT

c) $\mathcal{V} \circ \mathcal{N}$ is CP

↑

$$\mathcal{V}: X \mapsto X^T$$

Proof: a) \Rightarrow b) $(\text{id}_A \otimes \mathcal{N})(\rho_{AA'})$ is PPT, in particular for $\rho = \gamma_{AA'} \checkmark$

$$\text{b) } \Rightarrow \text{c) } (\text{id}_A \otimes \mathcal{V} \circ \mathcal{N})(\gamma) = (\text{id}_A \otimes \mathcal{V})(\text{id} \otimes \mathcal{N})(\gamma_{AA'})$$

$$= (\text{id}_A \otimes \mathcal{V})(\tau_{AB}^{\mathcal{N}})$$

$$= (\tau_{AB}^{\mathcal{N}})^T \geq 0 \quad \checkmark$$

$$\text{c) } \Rightarrow \text{a) } \mathcal{V} \circ \mathcal{N} \text{ is CP: } (\text{id}_R \otimes \mathcal{V} \circ \mathcal{N})(\rho_{RA}) \geq 0 \quad \forall \rho_{RA} \geq 0$$

$$\Rightarrow \mathcal{N} \text{ is PPT}$$

□

Remark: In general, $\mathcal{V} \circ \mathcal{N}$ is not CP.

However, for every CP map \mathcal{N} , the map $\mathcal{V} \circ \mathcal{N} \circ \mathcal{V}$ is CP. (Ex)

What about the capacities of PPT channels?

→ Horodecki³: PPT states are undistillable. $\nearrow \rho_{AB}^{\otimes n}$

Entanglement distillation: Given: iid copies of a state ρ_{AB}

Goal: convert $\rho_{AB}^{\otimes n}$ to a smaller number m_n of maximally entangled states, $|\phi^+\rangle = \frac{1}{\sqrt{2}}(|0\rangle|0\rangle + |1\rangle|1\rangle)$

If there is an entanglement distillation protocol (LOCC)

such that i) error of the protocol goes to 0 as $n \rightarrow \infty$,

ii) rate $c := \lim_{n \rightarrow \infty} \frac{m_n}{n} > 0$,

then ρ_{AB} is distillable.

Horodecki³
 \Rightarrow

$Q(N) = 0$ for PPT-channels N

(holds even if two-way CC is allowed!)

Remarks: → for quantum capacity $Q(N)$ we only allow forward / one-way CC.

→ If you also want to consider 2-way CC, then the relevant capacity is the two-way Q.cap. Q_2 (or Q_{\leftrightarrow})

→ $Q(N) \leq Q_2(N) \forall N$

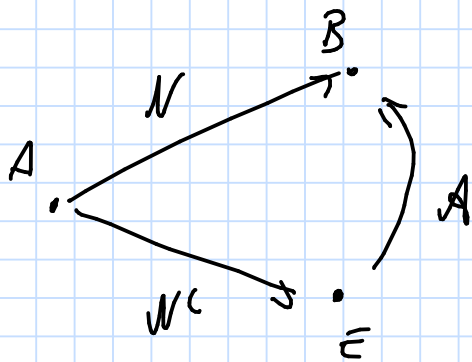
→ N is PPT $\Rightarrow Q(N) = Q_2(N) = 0$.

Classical capacity of PPT channels is generally unknown!

§2.5 Antidegradable channels

quantum channel $N: A \rightarrow B$ with an isometry $V: \mathcal{H}_A \rightarrow \mathcal{H}_B \otimes \mathcal{H}_E$
s.t. $N(\rho) = \text{tr}_E V \rho V^\dagger$.

Complementary channels: $N^c(\rho) = \text{tr}_B V \rho V^\dagger$



Def 13

N as above is called antidegradable,
if there is a channel $A: E \rightarrow B$ s.t.

$$N = A \circ N^c$$

Intuition: Eve (environment) can locally obtain Bob's output
via the channel A .

Antideg. channels cannot transmit quantum information: $Q(N) = 0$

"Proof": Assume $Q(N) > 0 \Leftrightarrow$ Alice can faithfully send qubits
to Bob at a positive rate. But there is a protocol
based on the channel A that lets Eve implement
the same protocol that Alice and Bob use.

This would clone Alice's state to B/E \Rightarrow violates no-cloning!
 \hookrightarrow "□"

Examples of anti-degradable channels:

- a) erasure channel $\mathcal{E}_p : \rho \mapsto (1-p)\rho + p \tau(\rho)$ for $p \geq \frac{1}{2}$.
- b) amplitude damping channel \mathcal{A}_γ for $\gamma \geq \frac{1}{2}$
- c) depolarizing channel \mathcal{D}_p for $p \geq \frac{1}{4}$.